On expansions of the group of integers

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A finite automaton M is a finite state machine given by (Σ, Q, q_0, F, T) , where Σ is a finite alphabet, q_0 an initial state, Q a finite set of states, $F \subset Q$ a set of accepting states and T a transition function from $Q \times \Sigma \rightarrow Q$ with the convention that $T(q, \lambda) = q$ for λ the empty word.

Let Σ^* be the set of all finite words on Σ . Let $\sigma \in \Sigma^*$ and $a \in \Sigma$, one extends T to $T : Q \times \Sigma^* \to Q$ by setting $T(q, \sigma a) := T(T(q, \sigma), a)$.

The automaton accepts $\sigma \in \Sigma^*$ if $T(q_0, \sigma) \in F$. (We will say that M works on Σ .)

A subset $L \subset \Sigma^*$ is recognized/accepted by M if $L = \{\sigma \in \Sigma^* : T(q_0, \sigma) \in F\}.$

Each natural number *n* can be written in base $d \ge 2$, so as a finite word in the alphabet $\Sigma := \{0, \ldots, d-1\}$. A subset of \mathbb{N} is *d*-automatic if it is recognized by a finite automaton *M* working on Σ .

Let d = 2. The study of 2-automatic subsets of \mathbb{N} has been marked by a result of R. Buchi who showed that the sets definable in $(\mathbb{N}, +, V_2)$, where $V_2(n)$ is the highest power of 2 dividing n, are exactly the 2-automatic sets.



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In fact the original statement of R. Buchi is about $(\mathbb{N}, +, P_2)$, where P_2 is the set of powers of 2, which is incorrect as observed, for instance, by Semenov.

Also, in his review, McNaughton suggested to consider instead the structure $(\mathbb{N}, +, \epsilon_2)$, where $\epsilon_2(n, m)$ holds iff n is a power of 2 which occurs in the binary expansion of m.

This last structure is interdefinable with $(\mathbb{N}, +, V_2)$.

A proof of the corrected result may be found in the master thesis of V. Bruyère.

The class \mathcal{R} of *regular* languages is the smallest class of languages on Σ that contains all finite languages and closed under the following operations: if $L_1, L_2 \in \mathcal{R}$, then

- $\ \, {\bf 0} \ \, L_1\cup L_2\in {\mathcal R},$

Fact: A subset $L \subset \Sigma^*$ is *regular* iff it is accepted by an automaton M working on Σ .

Using that correspondence, one can show that indeed a subset of $(\mathbb{N}, +, 0)$ is recognized by a finite automaton working on $\{0, 1\}$ iff it is definable in $(\mathbb{N}, +, 0, V_2)$. Furthermore that last structure is decidable. Also one can easily replace $(\mathbb{N}, +, 0)$ by $(\mathbb{Z}, +, -, 0, <)$.

Finite-automaton presentable structures, following B. Hodgson and A. Nies

Let \mathcal{L} be a finite language. A first-order (countable) \mathcal{L} -structure \mathcal{M} is *FA*-presentable (Finite Automaton) if there is a finite alphabet Σ and a regular language $D \subset \Sigma^*$ such that the elements of the domain can be represented by D and some finite automata can check whether the atomic relations hold in \mathcal{M} .

Theorem (Hodgson, 1976)

If a countable structure is FA-presentable, then it is decidable.

One uses that the emptiness problem for finite automata is decidable.

Remark: One can extend results on the ring \mathbb{Z} to other Euclidean rings, for instance polynomial rings or rings of formal power series over a finite field. (To treat that last example one needs to work with finite automata accepting infinite words).

Expansions of $(\mathbb{Z},+,<)$ and of $(\mathbb{R},+,\cdot,<)$

The theory of $(\mathbb{Z}, +, 0, <)$ satisfies different notions of minimality. Its theory is NIP, dp-minimal, coset-minimal and quasi-o-minimal. Recall that an ordered group *G* (non necessarily abelian) is coset-minimal if every definable subset of *G* is a finite union of cosets of definable subgroups intersected with intervals. A totally ordered structure is *quasi-o-minimal* if in every elementarily equivalent structure every definable set is a boolean combination of intervals and \emptyset -definable sets. It is *essentially quasi-o-minimal* if it has an expansion by constants which is quasi-o-minimal.

Fact: (Point-Wagner) Let G be an ordered group. Then the following is equivalent:

- Th(G) is essentially quasi-o-minimal.
- Th(G) is coset-minimal.
- Every definable subset of a model G of Th(G) is a finite union of cosets of nG intersected with intervals, for some $n < \omega$.

Note that $(\mathbb{Z}, +, <, V_2)$ has IP (one can define any finite subset of P_2 using ϵ_2).

 $(\mathbb{Z}, +, <, P_2)$ is NIP (see later).

(van den Dries, 1985) $Th(\mathbb{Z}, +, <, P_2)$ is model-complete in the language $\{+, -, <, \mod n, \lambda_2 : n \in \mathbb{N} \setminus \{0, 1\}\}$, where $\lambda_2(x)$ is the largest power of 2 smaller than x.

(Moosa-Scanlon, 2004) $Th(\mathbb{Z}, +, P_2)$ is stable and its definable subsets can be described as follows.

$(\mathbb{Z},+,P_2)$

Let $(\Gamma, +, -, 0)$ a finitely generated abelian group and F an injective endomorphism of Γ . Let Σ be a finite symmetric subset of Γ containing 0 and denote by Σ^* the set of finite words on Σ . We say that the finite word $\sigma := a_0 \cdots a_n$ represents $a \in \Gamma$ if $a = a_0 + F(a_1) + \ldots + F^n(a_n)$. We use the notation $a = [\sigma]_F$.

Definition (Bell-Moosa/Hawthorne)

Let Σ be as above. Then Σ is a *F*-spanning set for Γ if

- () any element of Γ can be represented by an element of $\Sigma^*,$
- \bigcirc if $a_1, a_2, a_3 \in \Sigma$, then $a_1 + a_2 + a_3 \in \Sigma + F(\Sigma)$,
- if $a_1, a_2 \in \Sigma$ are such that for some $b \in \Gamma$, $a_1 + a_2 = F(b)$, then $b \in \Sigma$.

PROPOSITION (Hawthorne)

The structure $(\Gamma, +, F)$ is FA-presentable.

Let $L \subset \Sigma^*$, then *L* is sparse if *L* is regular and if the set of words in *L* of length smaller than or equal to *x* is bounded by a polynomial function of *x*. A subset $A \subset \Gamma$ is *F*-sparse if $A = [L]_{F^r}$ for some sparse $L \subset \Sigma^*$ with Σ a F^r -spanning set, r > 0.

Let $a \in \Gamma$, then set $K(a, F) := \{a + F(a) + \cdots + F^n(a) \colon n \in \mathbb{N}\}.$

An elementary *F*-set is a subset of Γ of the form $a_0 + K(a_1, F^{n_1}) + \ldots + K(a_m, F^{n_m})$, for some $a_0, \ldots, a_m \in \Gamma$ and $n_1, \ldots, n_m \in \mathbb{N}$. Example: the *F*-orbit of *a* $F^{\mathbb{N}}(a) = \{F^n(a) : n \in \mathbb{N}\} = a + K(Fa - a)$ is an elementary *F*-set.

Theorem (Moosa-Scanlon 2004/Hawthorne 2021)

Let $A \subset \Gamma$ be *F*-sparse. Then *A* is stable in Γ iff *A* is a boolean combination of elementary *F*-sets.

As a consequence, one obtains:

Theorem (Hawthorne 2021–case d = 2)

Let $A \subset \mathbb{Z}$. Then A is 2-automatic and stable in $(\mathbb{Z}, +)$ iff A is a boolean combination of elementary 2-sets and cosets of subgroups of \mathbb{Z} iff A is definable in $(\mathbb{Z}, +, P_2)$. In particular the theory of $(\mathbb{Z}, +, P_2)$ is stable.

Independently, Poizat (2014) and then Palacin-Sklinos (2018) proved that:

Theorem (Poizat/Palacin-Sklinos)

The theory of $(\mathbb{Z}, +, P_2)$ is superstable of *U*-rank ω .

The result of L. van den Dries on the expansion $Th(\mathbb{Z}, +, \leq, P_2)$ was proven using a similar analysis as for $(\mathbb{R}, +, \cdot, 0, 1, \leq, P_2)$ (adding a multiplicative group to the field structure) which was extended later by van den Dries and Gunaydin to expansions of the form $(\mathbb{R}, +, \cdot, 0, 1, \leq, P_2.P_3)$.

Theorem (van den Dries, Gunaydin, 2006, Theorem 1.3)

Let Γ be a dense subgroup of $(\mathbb{R}^{>0})$ and suppose it has the Mann property. Then one can (explicitely) axiomatize the theory of $(\mathbb{R}, \Gamma) := (\mathbb{R}, +, -, \cdot, 0, 1, \Gamma, \{\gamma : \gamma \in \Gamma\}).$

Let K be a field and G a subgroup of the multiplicative group of K. Consider equations of the form:

$$r_1x_1+\ldots+r_nx_n=1,$$

with $r_i, 1 \le i \le n$, in the prime field of K and its solutions in G. A tuple $(g_1, ..., g_n) \in G^n$ is called a non-degenerate solution if for all proper subsets J of of $\{1, ..., n\}$, $\sum_{j \in J} r_j g_j \ne 0$. The group G has the Mann property in K if such equations have only finitely many non-degenerate solutions.

Fact [van der Poorten-Schlickewei, Evertse, Laurent] Any multiplicative group of finite rank in a field of characteristic 0 has the Mann property.

Theorem

Let K be an algebraically closed field of characteristic 0 and let Γ be a finitely generated subgroup of the K-points of some semi-abelian variety defined over K, then the theory of (K, Γ) is stable (Γ is stably embedded).

Theorem (Scanlon-Moosa, 2004)

Consider the structure $(K, +, \cdot, \Gamma)$, where K is an algebraically closed field of characteristic p and Γ is a finitely generated Frobenius submodule of a semi-abelian variety X defined over a finite field. Then the induced structure on Γ is stable and so the theory of (K, Γ) is stable.

How special is P_2 ?

Let $A \subset \mathbb{N}$, A infinite. Enumerate A as a strictly increasing sequence $A = (a_n)_{n \ge 0}$. Consider the expansions $(\mathbb{Z}, +, A)$, or $(\mathbb{Z}, +, -A \cup A)$, or $(\mathbb{Z}, +, 0, <, A)$. Which ones are tame?

How does the sequence grows? Consider the sequence $(\frac{a_{n+1}}{a_n})_{n\in\mathbb{N}}$ and $\limsup_{n\in\mathbb{N}} \frac{a_{n+1}}{a_n}$. If it is > 1, say that A is lacunary. If $\lim_{n\in\mathbb{N}} \frac{a_{n+1}}{a_n}$ exists in $\mathbb{R} \cup \{+\infty\}$, denote it by θ and call it the Kepler limit.

Recall that (a_n) is a linear recurrence sequence if there are $r_0, \ldots, r_{k-1} \in \mathbb{Q}$, with $k \in \mathbb{N}^{\geq 1}$ minimal, such that for all $n \in \mathbb{N}$,

$$a_{n+k}=\sum_{i=0}^{k-1}r_ia_{n+i}.$$

The polynomial P_A defined by $P_A(X) = X^k - \sum_{i=0}^{k-1} r_i X^i$ is called the *characteristic polynomial* and the elements a_0, \ldots, a_{k-1} the initial conditions.

A subset $B \subset \mathbb{R}^{>0}$ is geometric if $\{\frac{a}{b}: a \ge b \& a, b \in B\}$ is closed and discrete.

The sequence $A = (a_n) \subset \mathbb{Z}$ is a geometric sparse sequence if

• there is a function $f : A \to \mathbb{R}^{>0}$ such that f(A) is geometric, and

$$one sup \{ |a - f(a)| \colon a \in A \} \in \mathbb{R}.$$

Theorem (G. Conant, 2019)

Let A be a geometric sparse sequence (in \mathbb{Z}). Then, the theory of $(\mathbb{Z}, +, A)$ is superstable of Lascar rank ω .

Let $A := (a_n)_{n \in \mathbb{N}} \subset \mathbb{N}$. Then A is a regular (sparse) sequence if it has a Kepler limit exists and $\theta > 1$ and either

- θ is transcendental or
- θ is algebraic (over \mathbb{Q}) and (a_n) is a linear recurrence sequence whose characteristic polynomial is the minimal polynomial of θ

Theorem (Semenov(1979) /P.(2000); P.-Lambotte(2020))

Suppose A is a regular (sparse) sequence. Then, the theory of $(\mathbb{Z}, +, <, A)$ is model-complete and NIP, the theory of $(\mathbb{Z}, +, A)$ is superstable of Lascar rank ω and model-complete.

Decidability issues.

Sparse sequences following A.L. Semenov

Let S be the successor function on A and consider $Q(S(x)) := \sum_{i=0}^{n} r_i S^i(x)$ with $r_i \in \mathbb{Z}$, $0 \le i \le n$. We associate to such term a polynomial in $\mathbb{Z}[S]$.

Definition (A.L. Semenov)

 $A \subset \mathbb{N}$ is a sparse sequence if for any $Q(S) \in \mathbb{Z}[S] \setminus \{0\}$, either for all $a_m \in A$, $Q(a_m) = 0$, or $Q >_{pp} 0$ or $-Q >_{pp} 0$, and if $Q >_{pp} 0$, then there exists a natural number ℓ such that $Q(S^{\ell}) - Q > 0$.

One shows that if A is regular, then A is sparse following Semenov and that the theory T of $(\mathbb{Z}, +, -, 0, A, <)$ admits q.e. in the language of ordered abelian groups, together with $\{mod_n; n \in \mathbb{N}^*, \lambda_A, S, S^{-1}\}$. One can give an explicit axiomatisation and in case A is a linear recurrence sequence give conditions under which T is decidable. Again one can show that the theory of $(\mathbb{Z}, +, -, 0, A)$ admits q.e. in $\{+, -, 0, A\} \cup \{S, S^{-1}, c_0\} \cup \{\Sigma_{\bar{Q}} : \bar{Q} \in \mathbb{Z}[S]^n\}$, where $\Sigma_{\bar{Q}}$ are predicates that express that a tuple belongs to the images of $\bar{Q}(S)$.

(See the thesis of Q. Lambotte). There are regular sequences (a_n) which are not geometrically sparse. And among the linear recurrence sequences with a Kepler limit, one can characterize those which are geometrically sparse.

Theorem (Conant, 2018)

Let Γ be a finitely generated monoid of (\mathbb{N}, \cdot) and $A \subset \Gamma$ be infinite. Then the theory of $(\mathbb{Z}, +, A)$ is superstable of *U*-rank ω .

If there are $a, b \in \Gamma$ with $\log_a(b)$ is irrational, then Γ is non-lacunary (Furstenberg).

Non-lacunary sequences

Let $r \in \mathbb{R} \setminus \mathbb{Q}$, r > 1 and let $B_r := (\lfloor nr \rfloor)_{n \in \mathbb{N}^*}$ be a Beatty sequence.

Theorem (Gunaydin-Ozsahakyan, 2021)

 $(\mathbb{Z}, +, -, 0, 1, B_r)$ is NIP (unstable), decidable, admits q.e. adding countably many predicates $D_{n,1}, D_{n,0}, n \in \mathbb{N} \setminus \{0, 1\}$, and there is no intermediate structure betwen $(\mathbb{Z}, +)$ and $(\mathbb{Z}, +, B_r)$, where $D_{n,1}(x) := \exists y \in B_r \ ny = x$ and $D_{n,0}(x) := \exists y \notin B_r \ ny = x$.

They use the following tool. Let

$$e:\mathbb{Z}\to S^1:z\mapsto e^{rac{i2\pi z}{r}},$$

where S^1 is the circle of center (0,0) and radius 1 in \mathbb{R}^2 . Let C be the (anti-clockwise) cyclic order on S^1 and $Z_r := e(\mathbb{Z})$. Denote by (e(-1)1) the (cyclic) interval in S^1 .

Gunaydin and Ozsahakyan show that $(\mathbb{Z}, +, -, 0, 1, B_r)$ is isomorphic to $(e(\mathbb{Z}), \cdot, -1, 1, e(1), (e(-1)1) \cap e(\mathbb{Z}))$.

Let G be any dense subgroup of $(\mathbb{R}, +, 0, <)$, for instance take a subgroup generated by $\langle \mathbb{Z}, \mathbb{Z}/r \rangle$, where $r \in \mathbb{R} \setminus \mathbb{Q}$. Robinson and Zakon showed that the theory of (G, +, -, 0, <) can be axiomatized as an ordered abelian group which is *regularly dense*, namely for each positive integer n, x < y implies that there is an element z such that x < nz < y, and where for each $n \ge 2$, the indices [G : nG] (finite or infinite), are specified.

Theorem (Giraudet, Leloup, Lucas; Bélair, P.)

Let $\mathcal{G} := (G \cap [01[, +_1, <).$ Then $Th(\mathcal{G})$ is model-complete (in some extension of the language by infinitely many constants symbols), NIP and decidable if is r.e.

Non-lacunary sequences

Let \mathcal{P} be the set of prime numbers. Recall Dickson hypothesis (or the linear Schnitzel hypothesis):

(D) Let $f_1(x), \ldots, f_n(x)$ be irreducible linear polynomials over \mathbb{Z} , each having a positive leading coefficient. Suppose that there is no prime p which divides $\prod_{i=1}^{n} f_i(x)$ for all $x \in \mathbb{N}$. Then there exist infinitely many $x \in \mathbb{N}$ such that such that $f_1(x), \ldots, f_n(x)$ are all prime.

Theorem (Bateman-Jockusch-Woods, 1993)

 $(\mathbb{Z},+,0,<,\mathcal{P})$ is undecidable if Dickson hypothesis holds.

The proof shows that multiplication is definable, using a theorem of Buchi on the undecidability of $(\mathbb{Z}, +, <, 0, \{g(n) : n \in \mathbb{N}\})$, where $g(x) \in \mathbb{N}[x]$ is a polynomial of degree at least 2.

Theorem (Bateman-Jockusch-Woods, 1993)

 $(\mathbb{Z}, +, 0, <, V_2, \{2^n \colon n \in \mathcal{P})$ is decidable if Dickson hypothesis holds.

Theorem (Kaplan-Shelah, 2017)

 $(\mathbb{Z},+,0,1,-\mathcal{P}\cup\mathcal{P})$ is unstable, supersimple of rank 1 and decidable if Dickson hypothesis holds.

Note that in $(\mathbb{Z}, +, \mathcal{P})$ one may define the order (since any sufficiently large positive integer is a sum of at most 4 prime numbers).